Assessing trade-offs between energy consumption and security in sensor networks: simulations or testbeds?

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Laurent Mounier, Karine Altisen, Stéphane Devismes, Pascal Lafourcade, Raphaël Jamet, Ozgün Pinarer, ...

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Trade-off Between Security and Energy Consumption in Sensor Networks Applications...



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- I How to evaluate energy consumption of a complete HW/SW solution?
- 2 How to define and implement measurable security?

In such a way that we trust simulations/experiments.

We worked on points 1 and 2 in the ANR ARESA and ARESA2 projects, with Orange Labs; and also on point 1 in the ANR HELP project, with STMicroelectronics.

This Talk

 Brief presentation of one particular security protocol for WSNs, for which security has been evaluated on abstract models of WSNs

SR3: Secure Resilient Reputation-based Routing. Karine Altisen, Stéphane Devismes, Raphaël Jamet, Pascal Lafourcade - IEEE Int. Conf. on Distributed Computing in Sensor Systems (DCOSS 2013)

- More realistic assumptions and new questions on the attacker models
- Some (hopefully) more general comments/conclusions/questions

The SR³ Protocol

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Such a powerful attacker of a WSN will always win.



Weaker attackers considered:

"normal" nodes (w.r.t. HW capabilities) whose SW has been compromised Example: a node can start behaving as a blackhole at some point in time.

Expected properties of SR³: if there are not too many such nodes, the protocol will maintain a good delivery rate, at a reasonable cost.

The SR³ Protocol: Main Ideas



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Choice of a neighbor:

(i) part of the time: random; otherwise random weighted by reputation list.

The SR³ Protocol: Main Ideas



The SR³ Protocol: Simulations with Sinalgo

(A distributed algorithm simulator, ETH Zürich)



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A node that has acquired a good reputation turns into a blackhole: the delivery rate drops, and then SR3 recovers. Size of the reputation list: if increased, better routes, longer time to recover.

The SR³ Protocol: Energy Consumption

For the moment, the Sinalgo models are based on 2 important hypotheses:

- The radio channel and the MAC layer provide the routing level with an ideal platform (no collisions)
- Energy consumption is directly related to messages (number of messages sent, length of the routes)

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Next step: how would SR³ behave with more realistic assumptions:

- Realistic radio channel and MAC protocol
- $\bullet\,$ Energy consumption related to the behavior of the HW elements (CPU, radio, ...)

Experiments with SR^3

Sinalgo



WSNET (ideal MAC)



Experiments with SR³

Sinalgo



WSNET (ideal MAC)



More understandable

SensOrLabs testbed





Experiments with SR³

Sinalgo



WSNET (802.11)



SensOrLabs testbed



WSNET (ideal MAC)



+ "realistic" medium





Experiments with SR³



More understandable

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- Degraded performances due to the realistic radio channel + MAC *longer routes, ...*
- ② New behaviors due to interference effects hidden by the "ideal communication" hypothesis spatial and temporal correlations for the success rate of transmissions, implying weird behaviors of the reputation lists, ...

We try and build a situation for point 2 (topology, moment when a node becomes a blackhole, ...).

Unexpected Observation

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Very important impact of the attacker's position: — Where to put the attacker to make it as strong as possible?

An idea: We would like the honest nodes to suffer from the realistic channel, but NOT the attackers!

General Comments

The more details you can put in a reasonably fast simulation, the better?

No, because:

- More details means less understandable results
- More details in the execution environment (HW+channel+attackers) make it more particular, hence the SW is less robust/secure.

Sinalgo



WSNET (ideal MAC)



Sinalgo



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More understandable

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(drivers, low-level, MAC, routing, application)

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HW of a node radio channel

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Unfriendliness should be exagerated

(1) Exagerating Unfriendliness of the HW platform:

The very goal of "transaction-level modeling" (TLM) for systems-on-a-chip.
 Feasible even though the REAL object code runs on rather abstract HW models (bit-accurate models)

Example with two timers:

Embedded SW:

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start timer 1 delay 10
start timer 2 delay 20
wait expiration 1
wait expiration 2
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HW models of increasing unfriendliness:

- 1) exact timing: 10, 20
- 2) loose timing: [8,12], [18,22]
- 3) no timing: one day, one day

If even loose timing is considered too friendly, then untimed models are the only choice to guarantee SW robustness.

(2) Exagerating unfriendliness of the attackers in a WSN?

We could imagine a mixed simulation mode in which:

— The honest nodes play by the rules, i.e. they suffer from the imperfections of the radio channel+MAC

- The attackers don't play by the rule, i.e., for them the transmission is perfect.

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For the attackers: if any position with the realistic radio channel is too particular/friendly, the only choice is to build an *ideal* position .

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A very useful tool would be a mix of:

- testbeds and
- high level models with perturbations
 - $= \mbox{exagerated}$ view of the discrepancies between ideal and real behaviors

Questions?